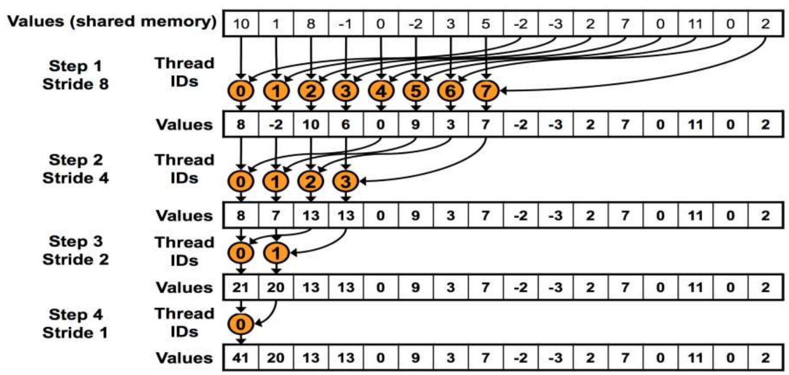
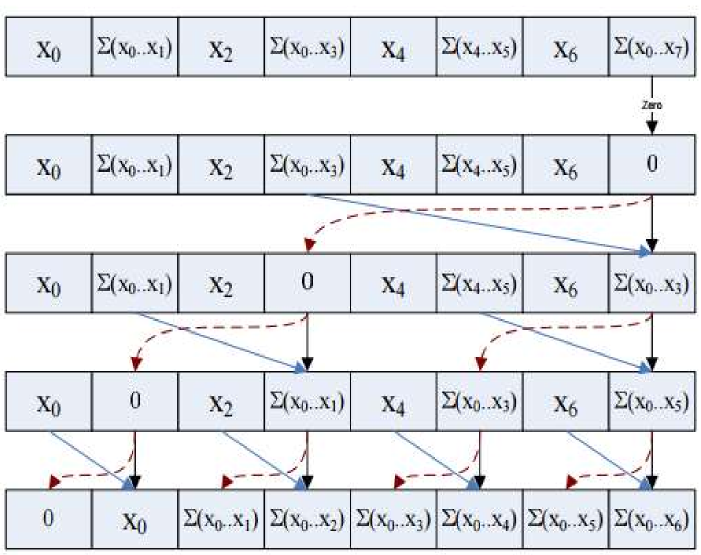
CS 179: GPU Programming

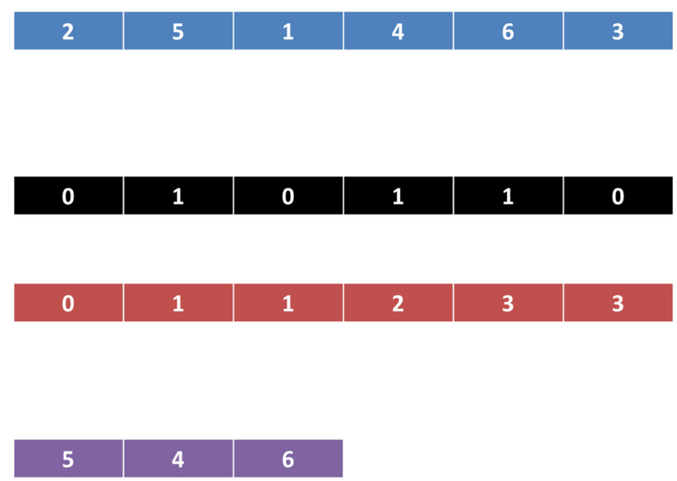
Lecture 8

Last time





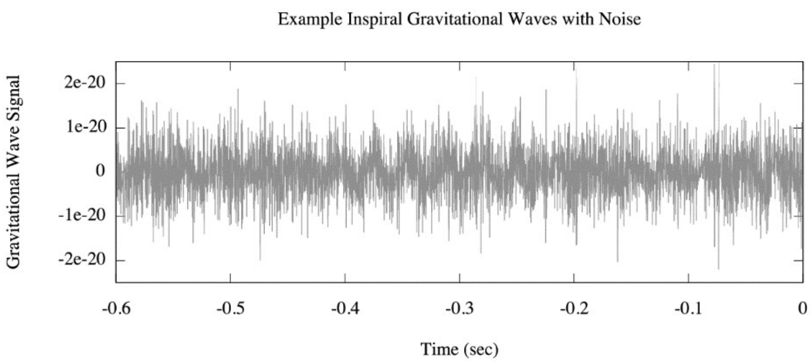
- GPU-accelerated:
  - Reduction
  - Prefix sum
  - Stream compaction
  - Sorting (quicksort)

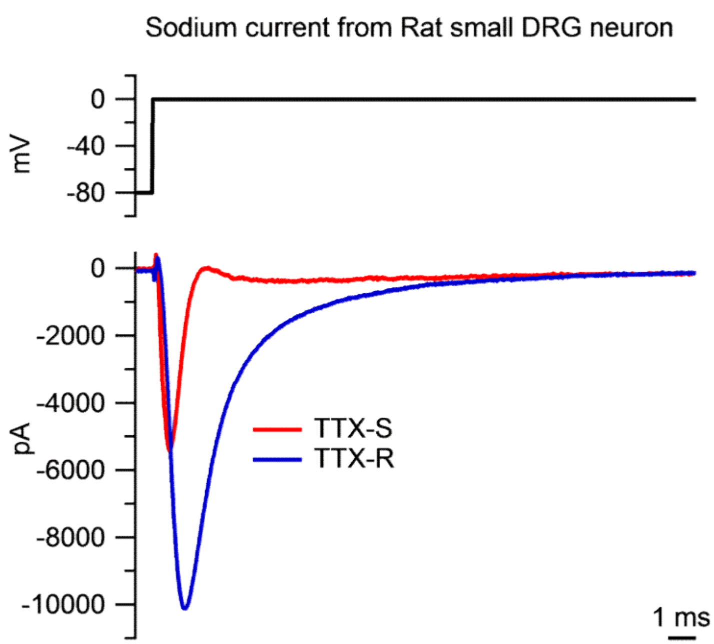


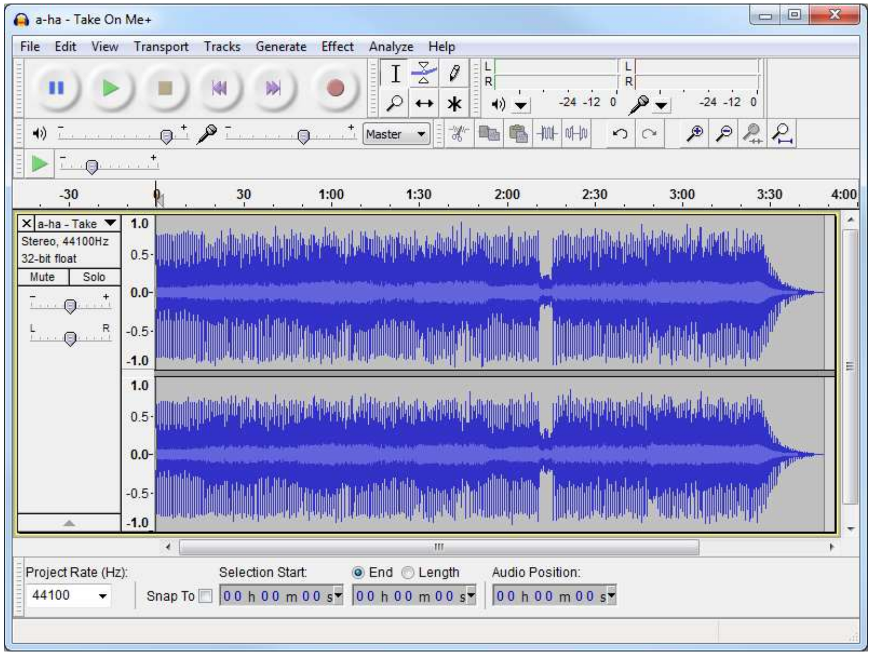
Today

- GPU-accelerated Fast Fourier Transform
- cuFFT (FFT library)

Signals (again)







“Frequency content”

- What frequencies are present in our signals?

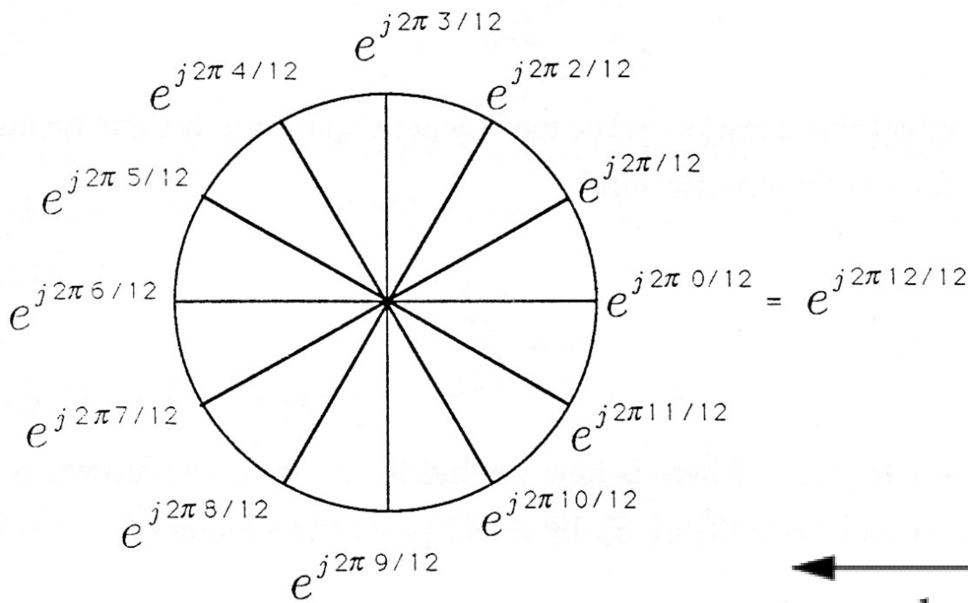
Discrete Fourier Transform (DFT)

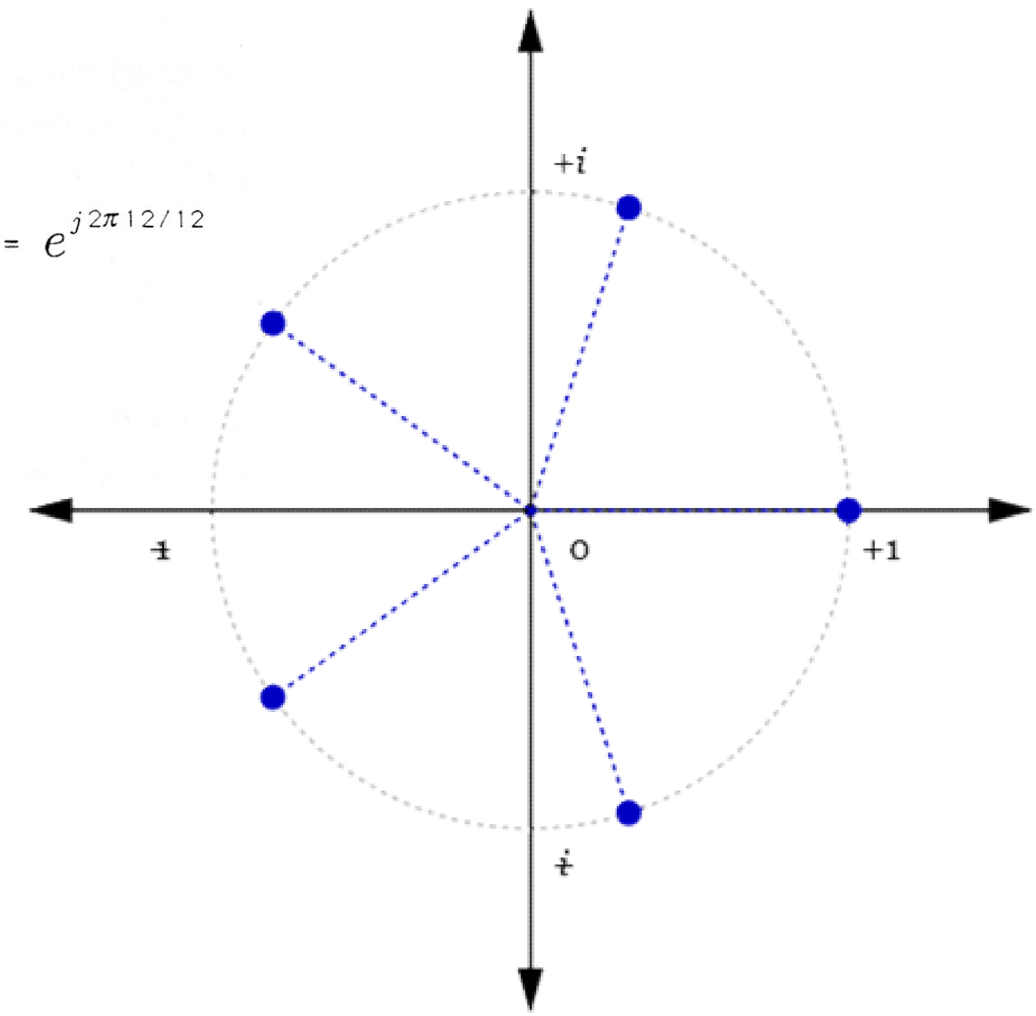
$$W = \frac{1}{\sqrt{N}} \begin{bmatrix} 1 & 1 & 1 & 1 & \cdots & 1 \\ 1 & \omega & \omega^2 & \omega^3 & \cdots & \omega^{N-1} \\ 1 & \omega^2 & \omega^4 & \omega^6 & \cdots & \omega^{2(N-1)} \\ 1 & \omega^3 & \omega^6 & \omega^9 & \cdots & \omega^{3(N-1)} \\ \vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & \omega^{N-1} & \omega^{2(N-1)} & \omega^{3(N-1)} & \cdots & \omega^{(N-1)(N-1)} \end{bmatrix},$$

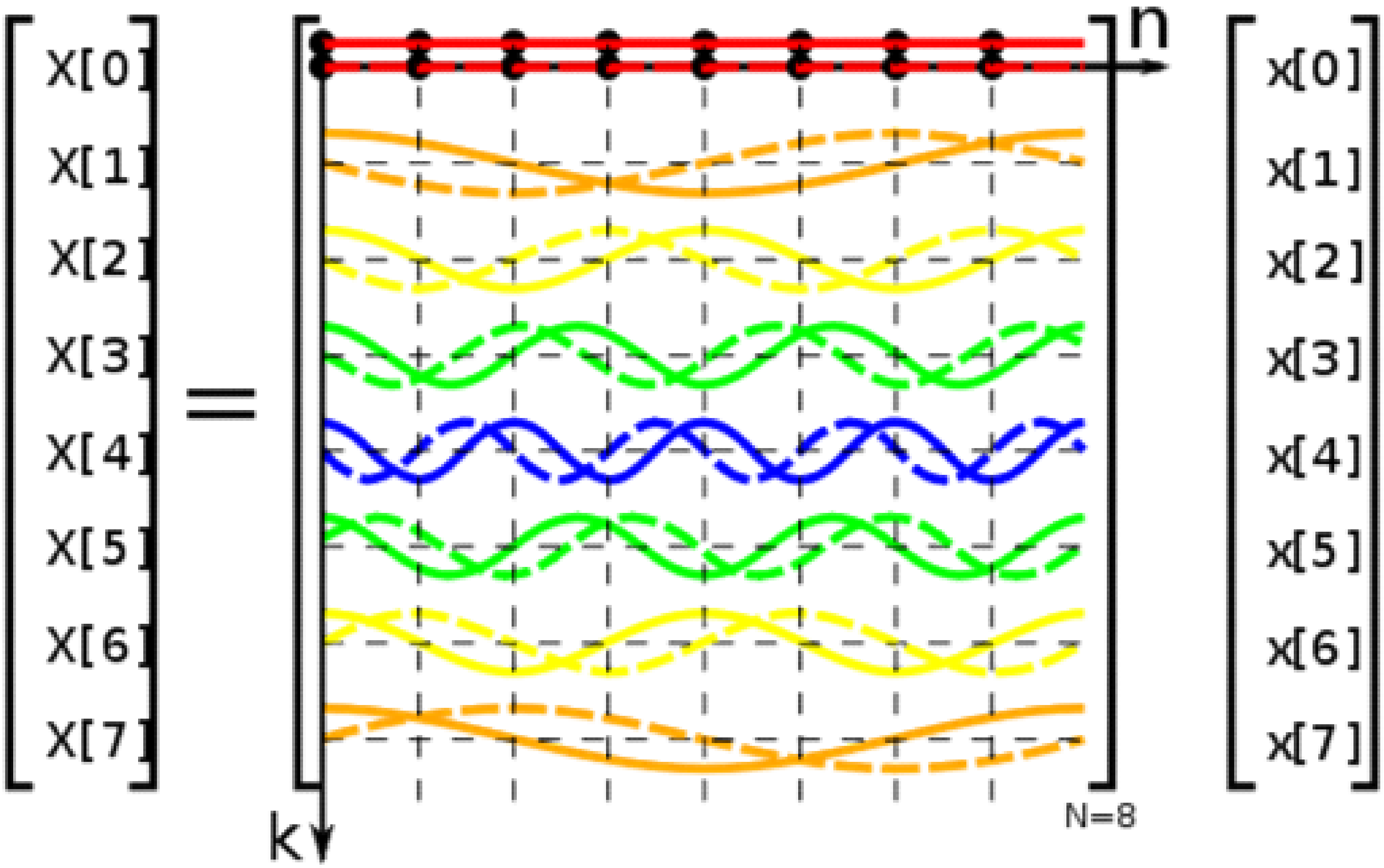
$\omega = e^{-2\pi \ /N}$

- Given signal $\vec{x} = (x_1, \dots, x_N)$ over time, $\vec{y} = W\vec{x}$ represents DFT of $\vec{x}$
  - Each row of *W* is a complex sine wave
  - Each row multiplied with $\vec{x}$ - inner product of wave with signal
  - Corresponding entries of $\vec{y}$ - “content” of that sine wave!

Roots of unity






X[0]
X[1]
X[2]
X[3]
X[4]
X[5]
X[6]
X[7]
=
n
k
N=8
x[0]
x[1]
x[2]
x[3]
x[4]
x[5]
x[6]
x[7]

Discrete Fourier Transform (DFT)

- Alternative formulation:

$$X_k \stackrel{\text{def}}{=} \sum_{n=0}^{N-1} x_n \cdot e^{-2\pi i k n / N}, \quad k \in \mathbb{Z}$$

  - X_k - values corresponding to wave *k*
    - Periodic – calculate for $0 \leq k \leq N - 1$

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  – Naive runtime: $O(N^2)$
    - Sum of *N* iterations, for *N* values of *k*

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  - X_k - values corresponding to wave *k*
    - Periodic – calculate for 0 ≤ *k* ≤ *N* - 1

Number of distinct values: N, not N^2 !

(Proof)

- Breakdown (assuming N is power of 2):
  - (Let $\omega_N = e^{-2\pi i/N}$, smallest root of unity)

$$\sum_{n=0}^{N-1} x_n {\omega_N}^{kn}$$

(Proof)

- Breakdown (assuming N is power of 2):
  - (Let $\omega_N = e^{-2\pi i/N}$, smallest root of unity)

$$\sum_{n=0}^{N-1} x_n {\omega_N}^{kn}$$

$$= \sum_{n=0}^{N/2-1} x_{(2n)} {\omega_N}^{k(2n)} + \sum_{n=0}^{N/2-1} x_{(2n+1)} {\omega_N}^{k(2n+1)}$$

(Proof)

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$$= \sum_{n=0}^{N/2-1} x_{(2n)} {\omega_N}^{k(2n)} + \omega_N \sum_{n=0}^{N/2-1} x_{(2n+1)} {\omega_N}^{k(2n)}$$

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$$= \sum_{n=0}^{N/2-1} x_{(2n)} {\omega_{N/2}}^{kn} + \omega_N \sum_{n=0}^{N/2-1} x_{(2n+1)} {\omega_{N/2}}^{kn}$$

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$$= \sum_{n=0}^{N/2-1} x_{(2n)} {\omega_N}^{k(2n)} + \omega_N \sum_{n=0}^{N/2-1} x_{(2n+1)} {\omega_N}^{k(2n)}$$

$$= \underbrace{\sum_{n=0}^{N/2-1} x_{(2n)} {\omega_{N/2}}^{kn}}_{\text{DFT of } x_n \text{, even n!}} + \omega_N \underbrace{\sum_{n=0}^{N/2-1} x_{(2n+1)} {\omega_{N/2}}^{kn}}_{\text{DFT of } x_n \text{, odd n!}}$$

(Divide-and-conquer algorithm)

```
Recursive-FFT(Vector x):

    if x is length 1:
         return x

    x_even <- (x0, x2, ..., x_(n-2) )
    x_odd <- (x1, x3, ..., x_(n-1) )

    y_even <- Recursive-FFT(x_even)
    y_odd <- Recursive-FFT(x_odd)

    for k = 0, …, (n/2)-1:
        y[k]        <- y_even[k] + w^k * y_odd[k]
        y[k + n/2]  <- y_even[k] - w^k * y_odd[k]

    return y
```

(Divide-and-conquer algorithm)

```
Recursive-FFT(Vector x):

    if x is length 1:
         return x

    x_even <- (x0, x2, ..., x_(n-2) )
    x_odd <- (x1, x3, ..., x_(n-1) )

    y_even <- Recursive-FFT(x_even)             T(n/2)
    y_odd <- Recursive-FFT(x_odd)               T(n/2)

    for k = 0, …, (n/2)-1:                      O(n)
        y[k]         <- y_even[k] + w^k * y_odd[k]
        y[k + n/2]   <- y_even[k] - w^k * y_odd[k]

    return y
```

Runtime

- Recurrence relation:
  - $T(n) = 2T(n/2) + O(n)$

$O(n \log n)$ runtime! *Much* better than $O(n^2)$

- (Minor caveat: N must be power of 2)
  - Usually resolvable

Parallelizable?

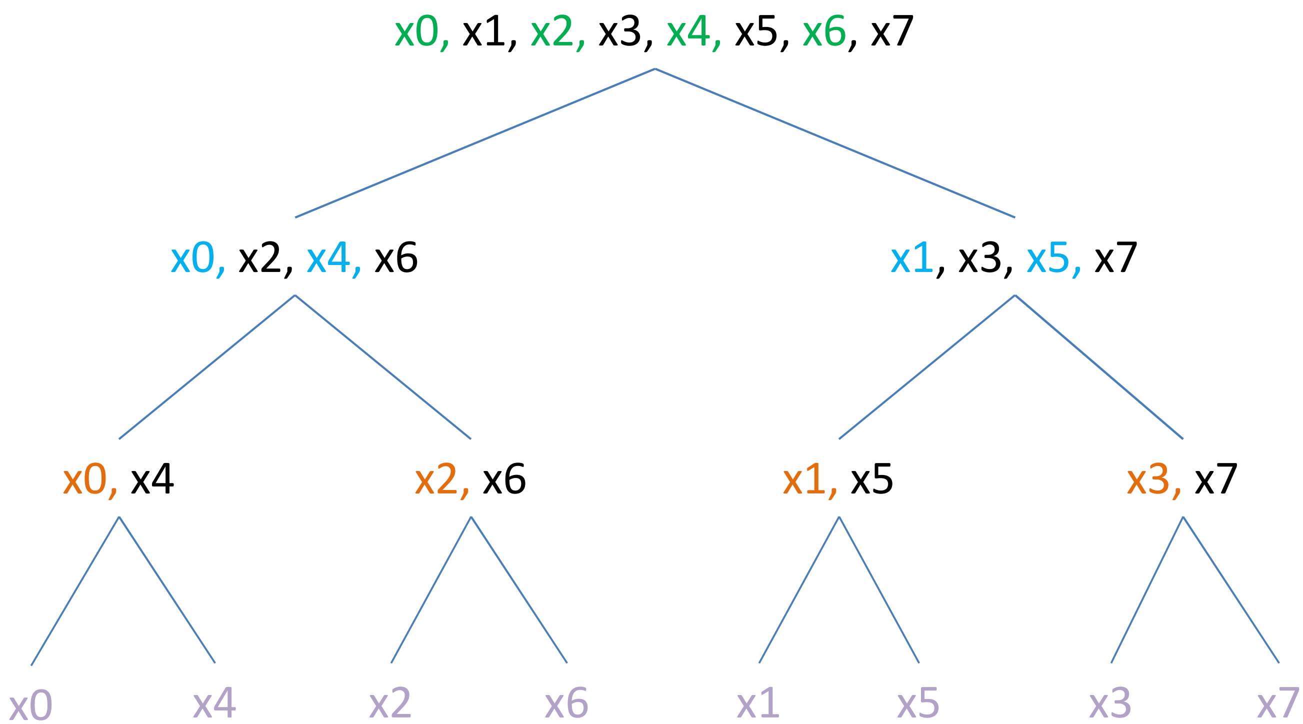
- $O(n^2)$ algorithm certainly is!

```
for k = 0,…,N-1:
    for n = 0,…,N-1:
        ...
```

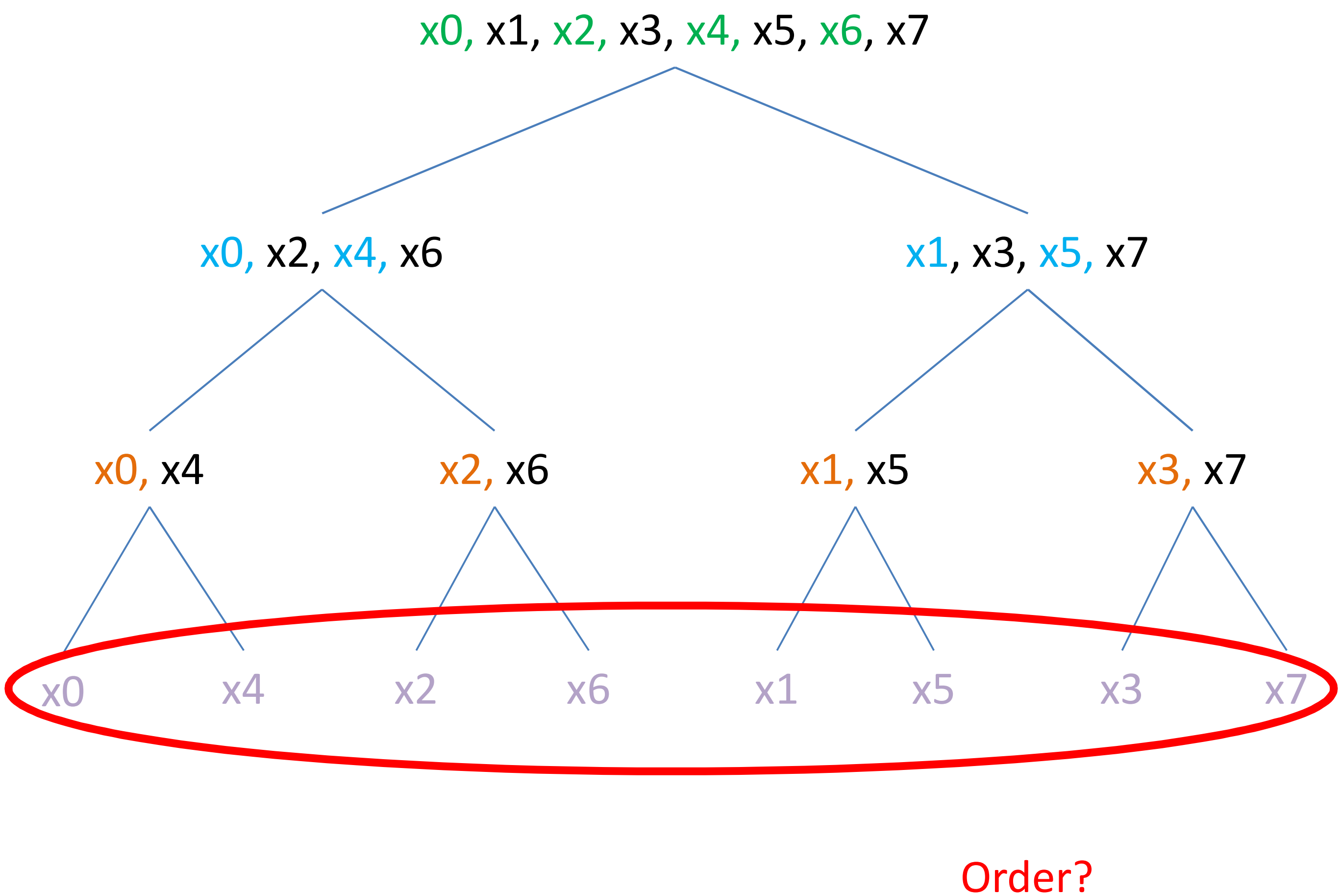
$$X_k \stackrel{\text{def}}{=} \sum_{n=0}^{N-1} x_n \cdot e^{-2\pi ikn/N}$$

- Sometimes parallelization outweighs runtime!
  – (N-body problem, …)

Recursive index tree



Recursive index tree

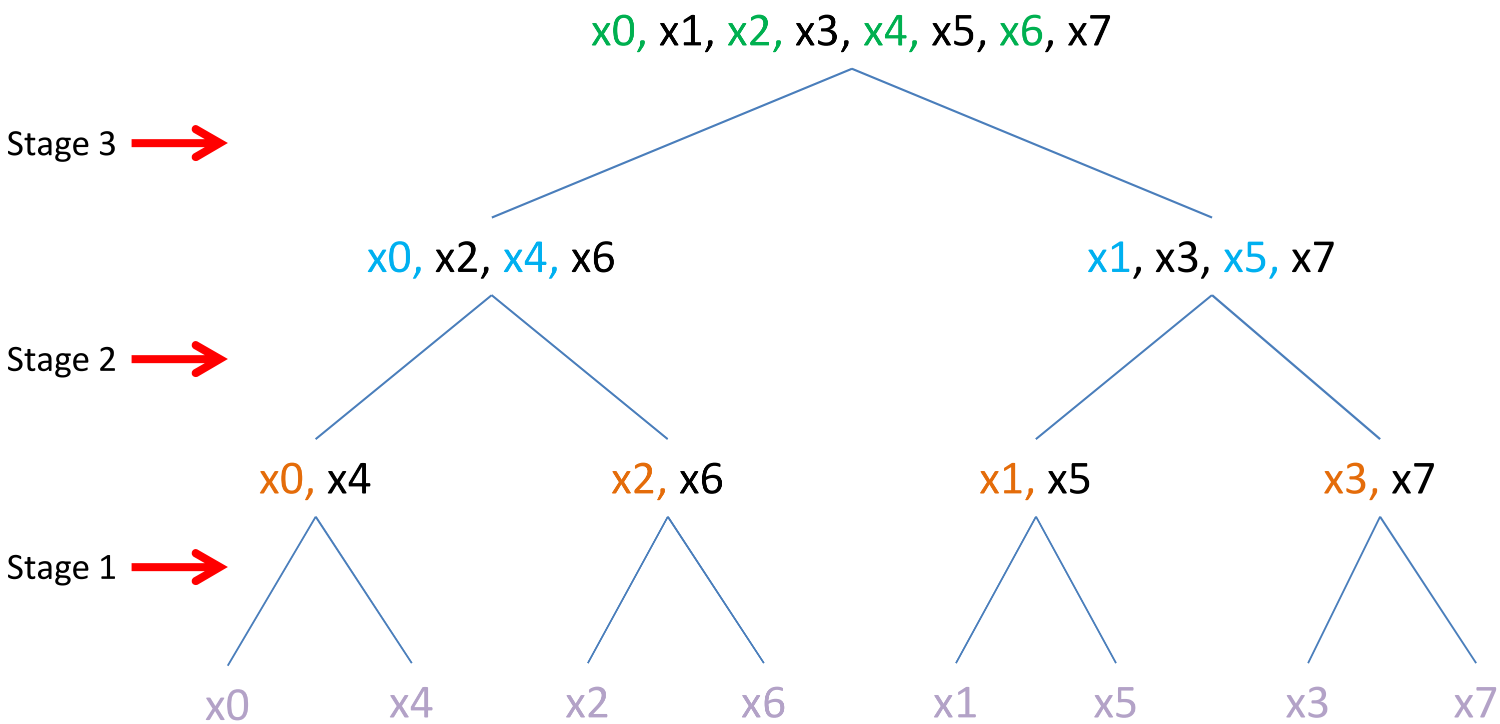


0	000
4	100
2	010
6	110
1	001
5	101
3	011
7	111

Bit-reversal order

0	000	reverse of…	000
4	100		001
2	010		010
6	110		011
1	001		100
5	101		101
3	011		110
7	111		111

Iterative approach



(Divide-and-conquer algorithm review)

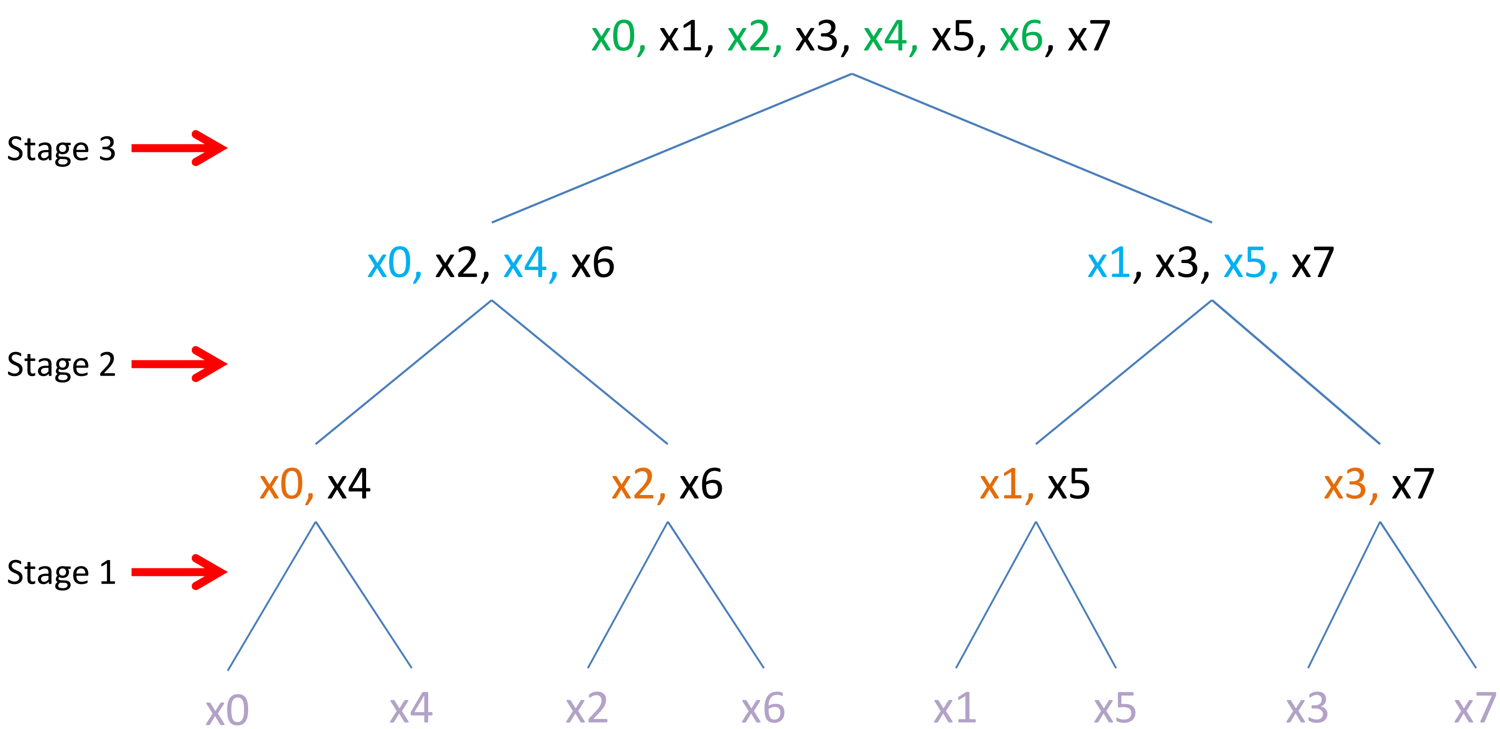
```
Recursive-FFT(Vector x):

    if x is length 1:
         return x

    x_even <- (x0, x2, ..., x_(n-2) )
    x_odd <- (x1, x3, ..., x_(n-1) )

    y_even <- Recursive-FFT(x_even)
    y_odd <- Recursive-FFT(x_odd)

    for k = 0, …, (n/2)-1:
        y[k]         <- y_even[k] + w^k * y_odd[k]
        y[k + n/2]   <- y_even[k] - w^k * y_odd[k]

    return y
```

T(n/2)

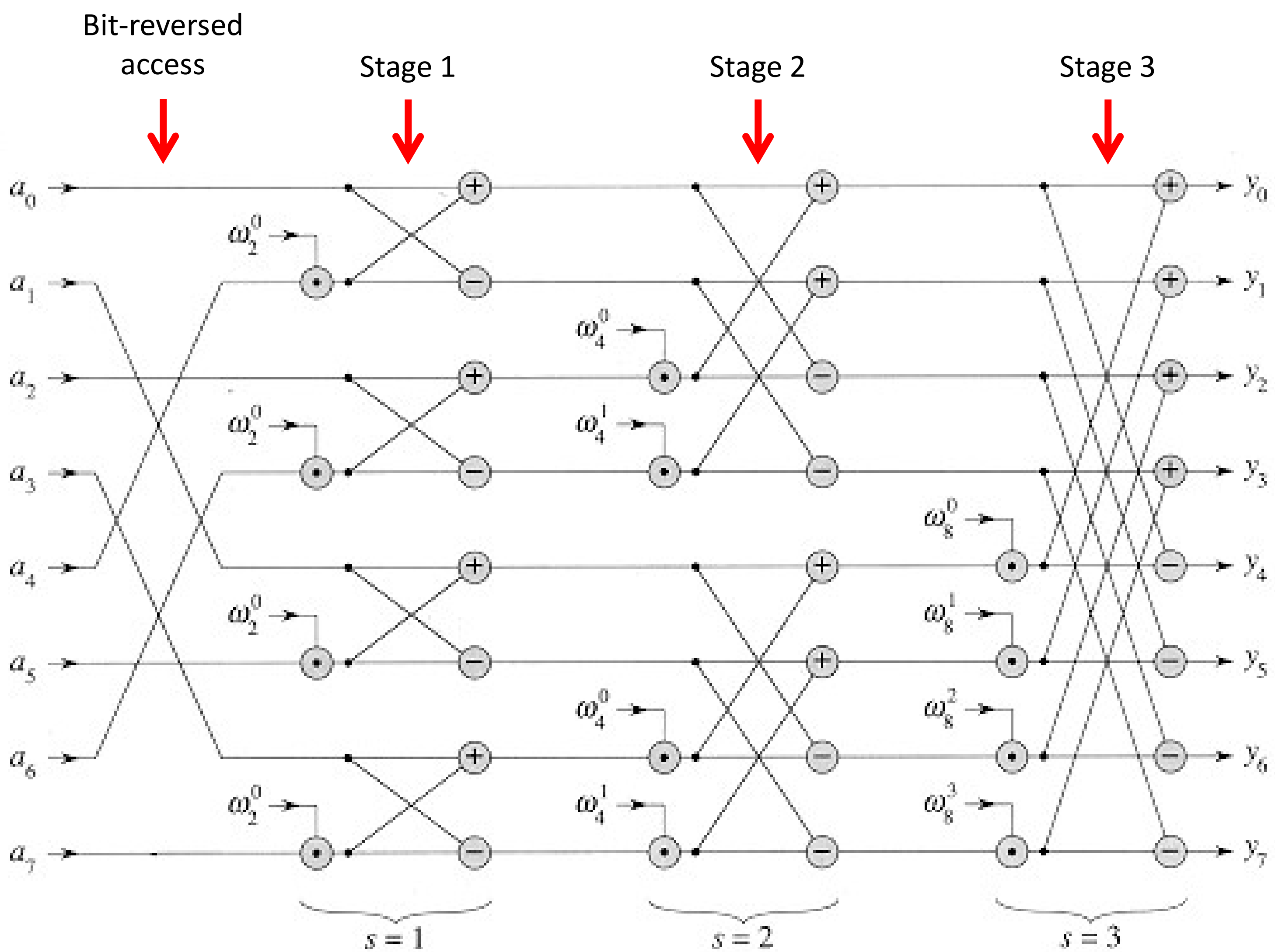
T(n/2)

O(n)

Iterative approach

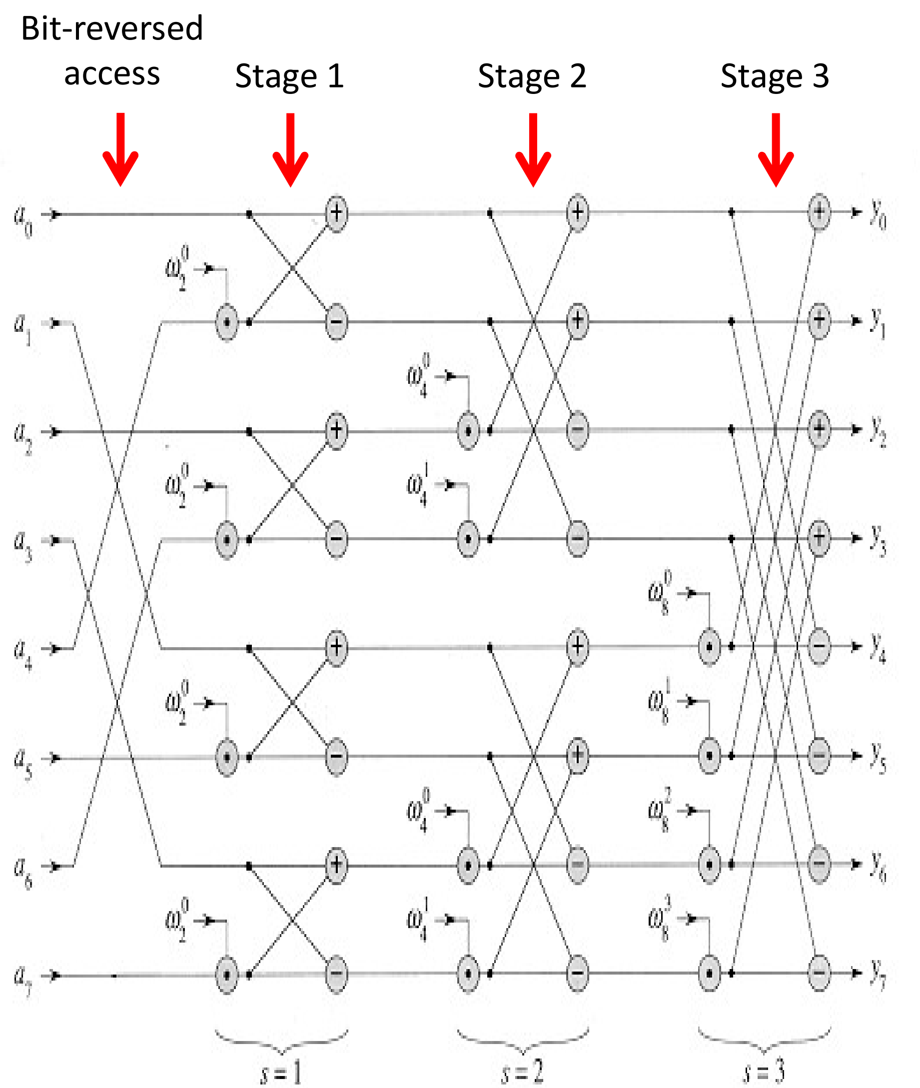


Iterative approach



http://staff.ustc.edu.cn/~csli/graduate/algorithms/book6/chap32.htm

Iterative approach



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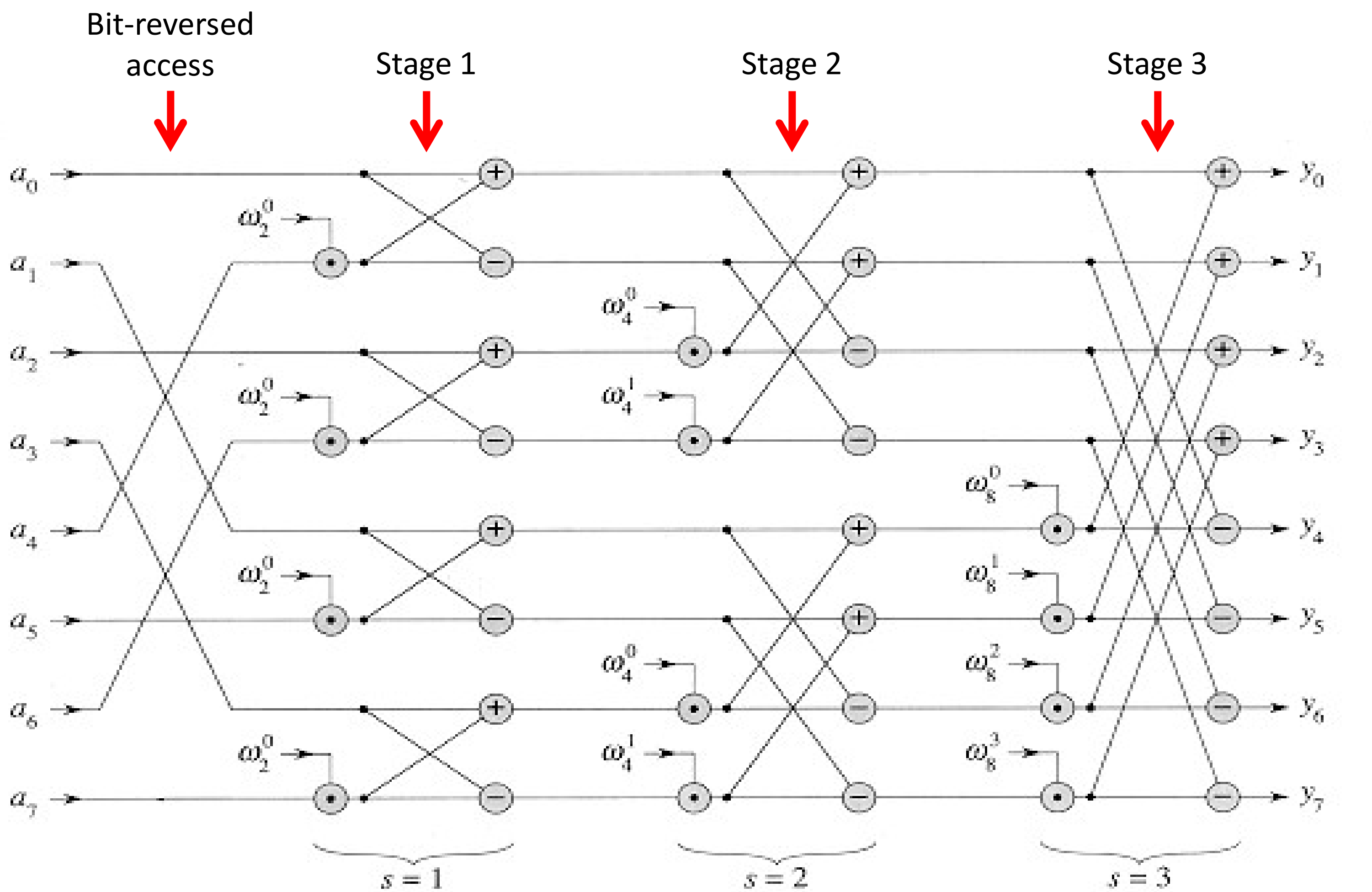
```
Iterative-FFT(Vector x):

    y <- (bit-reversed order x)
    N <- y.length
    for s = 1,2,…,log(N):

        m <- 2^s
        w_n <- e^(2πj/m)

        for k: 0 ≤ k ≤ N-1, stride m:
            for j = 0,…,(m/2)-1:

                u <- y[k + j]
                t <- (w_n)^j * y[k + j + m/2]

                y[k + j] <- u + t
                y[k + j + m/2] <- u - t


    return y
```

CUDA approach



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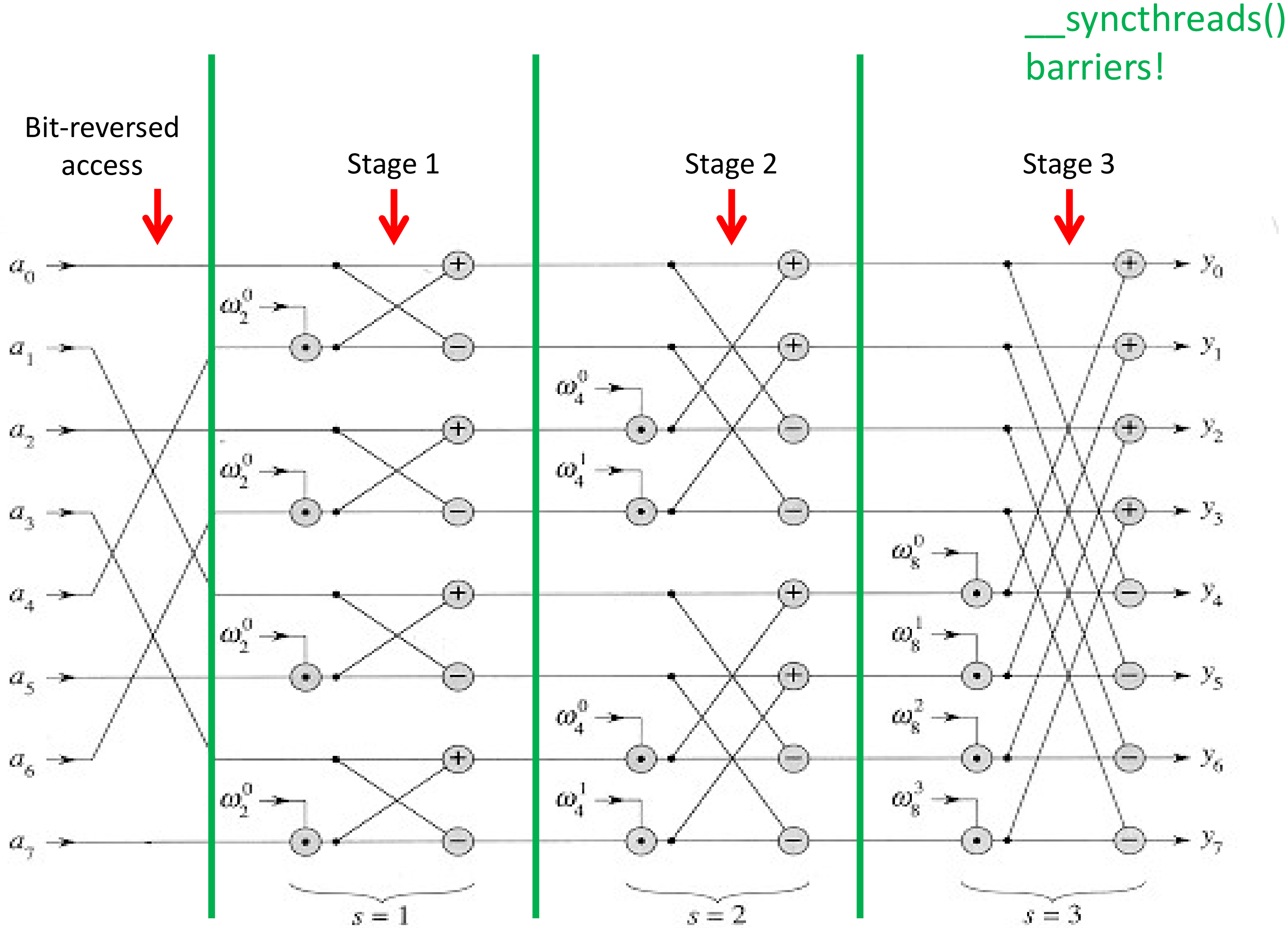
CUDA approach



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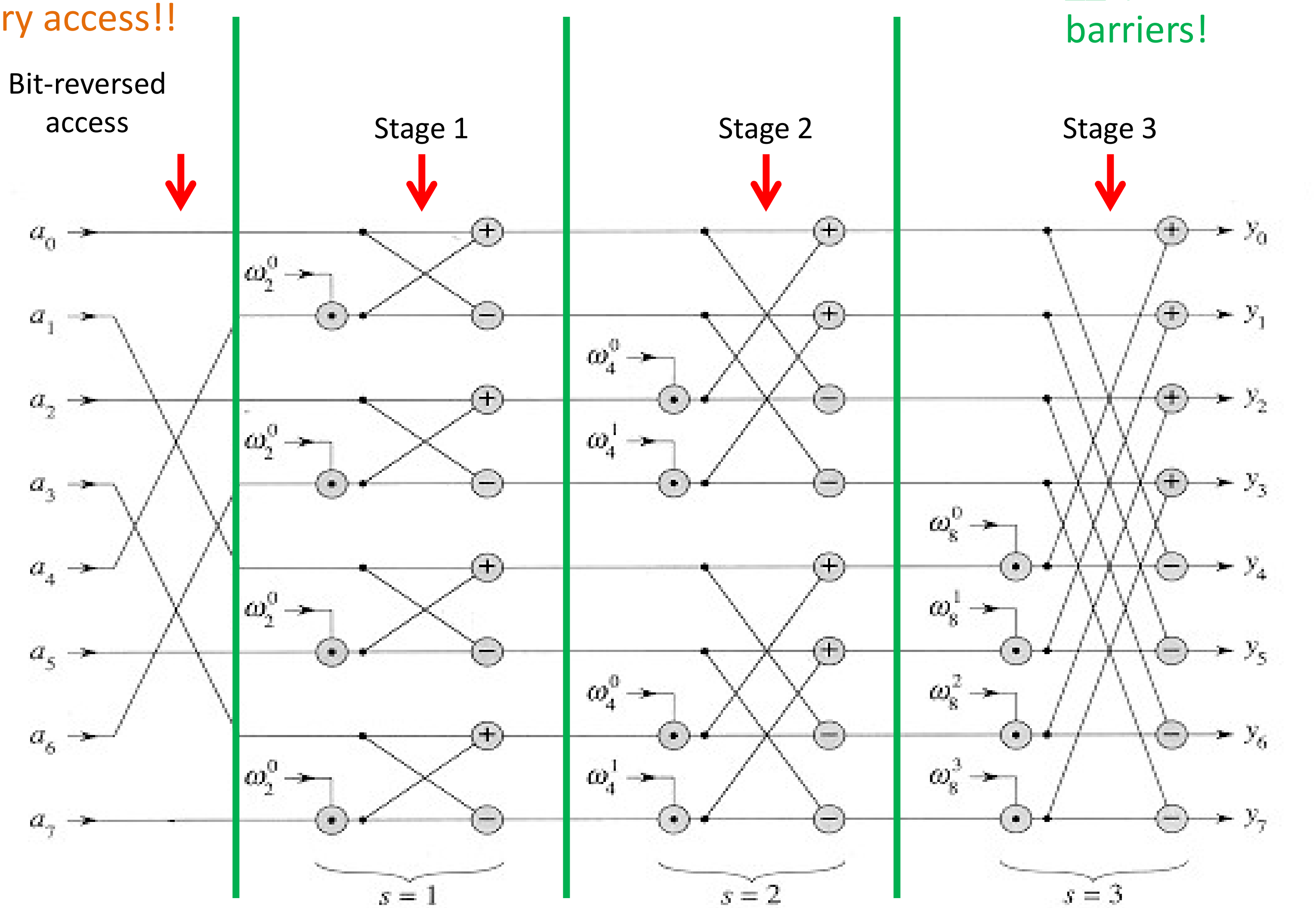
CUDA approach

Non-coalesced memory access!!

__syncthreads() barriers!



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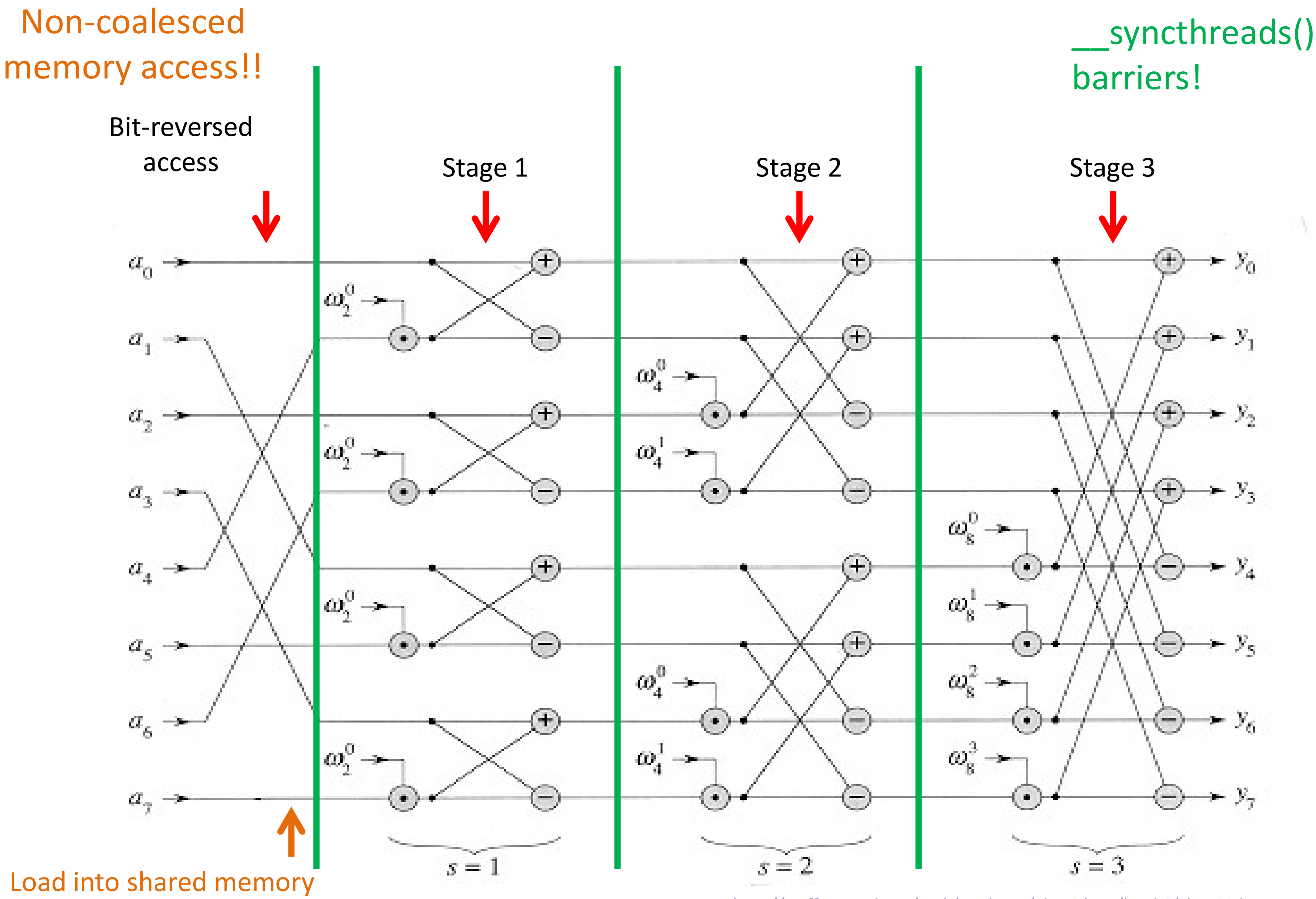
CUDA approach



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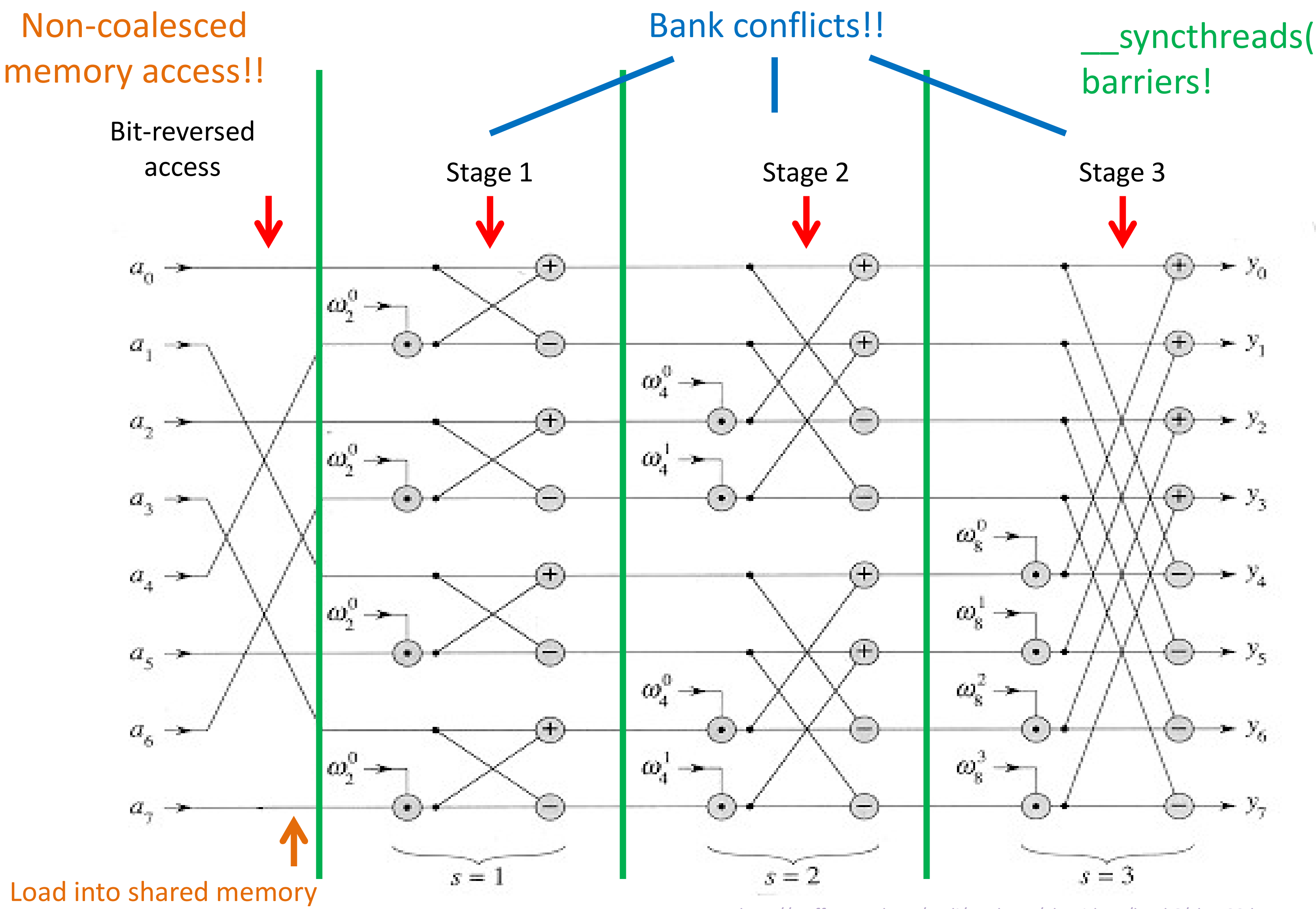

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Inverse DFT/FFT

- Similarly parallelizable!
  - (Sign change in complex terms)

$$x_n = \frac{1}{N} \sum_{k=0}^{N-1} X_k \cdot e^{i2\pi kn/N}, \quad n \in \mathbb{Z}$$

cuFFT

- FFT library included with CUDA
  - Approximately implements previous algorithms
    - (Cooley-Tukey/Bluestein)
    - Also handles higher dimensions

cuFFT 1D example

```
#define NX 262144

cufftComplex *data_host
        = (cufftComplex*)malloc(sizeof(cufftComplex)*NX);
cufftComplex *data_back
        = (cufftComplex*)malloc(sizeof(cufftComplex)*NX);


// Get data...


cufftHandle plan;
cufftComplex *data1;
cudaMalloc((void**)&data1, sizeof(cufftComplex)*NX);
cudaMemcpy(data1, data_host, NX*sizeof(cufftComplex), cudaMemcpyHostToDevice);

/* Create a 1D FFT plan. */
int batch = 1;  // Number of transforms to run
cufftPlan1d(&plan, NX, CUFFT_C2C, batch);

/* Transform the first signal in place. */
cufftExecC2C(plan, data1, data1, CUFFT_FORWARD);


/* Inverse transform in place. */
cufftExecC2C(plan, data1, data1, CUFFT_INVERSE);

cudaMemcpy(data_back, data1, NX*sizeof(cufftComplex), cudaMemcpyDeviceToHost);
```

Correction:
Remember to use
cufftDestroy(plan)
when finished with
transforms

cuFFT 3D example

```
#define NX 64
#define NY 64
#define NZ 128

cufftComplex *data_host
        = (cufftComplex*)malloc(sizeof(cufftComplex)*NX*NY*NZ);
cufftComplex *data_back
        = (cufftComplex*)malloc(sizeof(cufftComplex)*NX*NY*NZ);


// Get data...


cufftHandle plan;
cufftComplex *data1;
cudaMalloc((void**)&data1, sizeof(cufftComplex)*NX*NY*NZ);
cudaMemcpy(data1, data_host, NX*NY*NZ*sizeof(cufftComplex), cudaMemcpyHostToDevice);

/* Create a 3D FFT plan. */
cufftPlan3d(&plan, NX, NY, NZ, CUFFT_C2C);

/* Transform the first signal in place. */
cufftExecC2C(plan, data1, data1, CUFFT_FORWARD);


/* Inverse transform in place. */
cufftExecC2C(plan, data1, data1, CUFFT_INVERSE);

cudaMemcpy(data_back, data1, NX*NY*NZ*sizeof(cufftComplex), cudaMemcpyDeviceToHost);
```

Correction:
Remember to use
cufftDestroy(plan)
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Remarks

- As before, some parallelizable algorithms don’t easily “fit the mold”
  - Hardware matters more!
  - Some resources:
    - Introduction to Algorithms (Cormen, et al), aka “CLRS”, esp. Sec 30.5
    - “An Efficient Implementation of Double Precision 1-D FFT for GPUs Using CUDA” (Liu, et al.)