## CS 153 Current topics in theoretical computer science

Solution Set 1

Out: May 15

- 1. We use the substitution method. By setting  $b_{1,1} = 1$  and  $b_{2,1} = 0$ , we see that some product of *a*'s and *b*'s depends on  $a_{1,1}$ ; set  $a_{1,1}$  to a linear form in the *a*'s that makes this product zero. Then by setting  $b_{1,1} = 0$  and  $b_{2,1} = 1$ , we see that some product of *a*'s and *b*'s depends on  $a_{1,2}$ ; set  $a_{1,2}$  to a linear form in the *a*'s that makes this product zero. What remains still computes a  $1 \times 2$  by  $2 \times 2$  matrix multiplication. Since  $\langle 1, 2, 2 \rangle$  has 4 linearly independent slices, there must be at least 4 remaining slices.
- 2. Set  $\alpha_i = \log_{|G_i|}(|X_i| \cdot |Y_i| \cdot |Z_i|)$ . From a theorem proved in class, we have that for each i,

$$(|X_i||Y_i||Z_i|)^{\omega/3} \le D_i^{\omega-2}|G_i|.$$

Taking logs and dividing by  $\log |G_i|$  we get

$$\alpha_i \omega/3 \le (1/2 - \epsilon)(\omega - 2) + 1$$

Replacing  $\alpha_i$  with its supremum, we get

$$\omega/2 \le (1/2 - \epsilon)(\omega - 2) + 1$$

which simplifies to  $\epsilon \omega \leq 2\epsilon$ , hence  $\omega = 2$ .

3. Recall that we found three subsets X, Y, Z of  $S_n$  with  $|X| = |Y| = |Z| = |S_n|^{1/2-o(1)}$ . For each  $y \in Y$ , define:

$$A_y = \{xy^{-1} : x \in X\} B_y = \{yz^{-1} : z \in Z\}$$

Observe that  $A_y B_y = \{xz^{-1} : x \in X, z \in Z\}$ , and these must all be distinct; if we had  $(x, z) \neq (x', z')$  with  $xz^{-1} = x'(z')^{-1}$  then  $(x')^{-1}xz^{-1}z' = 1$  which violates the triple product property (since  $1 \in Q(Y)$ ). Also, if  $A_u B_u \cap A_y B_{y'}$  with  $y \neq y'$ , then we have  $xz^{-1} = x'y^{-1}y'(z')^{-1}$ , and thus  $x^{-1}x'y^{-1}y'(z')^{-1}z = 1$ , which violates the triple product property.

4. Such a table has three types of columns – columns containing only 1's and 2's, columns containing only 2's and 3's, and columns containing only 3's and 1's. Let  $n_1, n_2, n_3$  denote the number of each type of column. The number of distinct 1/2 patterns is  $2^{n_1}$ , the number of distinct 2/3 patterns is  $2^{n_2}$  and the number of distinct 3/1 patterns is  $2^{n_3}$ . If  $N > 2^{n_1}2^{n_2}$  then by the pigeonhole principle, there are two rows with identical 1/2 patterns in the 1/2 columns and 2/3 patterns in the 2/3 columns. Thus these two rows have identical "2-sets", and thus the table is not a strong USP. Thus  $N \leq 2^{n_1+n_2}$ . Similarly, the fact that a USP can have no duplicate "1-sets" implies  $N \leq 2^{n_1+n_3}$ , and the fact that a USP can have no duplicate "2-sets" implies that  $N \leq 2^{n_2+n_3}$ . Thus  $N^3 \leq 2^{2(n_1+n_2+n_3)} = 2^{2n}$ .

- 5. (a) Consider the matrix M with  $M[i, j] = \omega^{i+j}$ , where  $\omega$  is a primitive *n*-th root of unity, and note that M has rank 1. Let J be the all-ones matrix (also rank 1). Then M I has rank at most 2, and it has the same support as J I, which has rank n 1.
  - (b) We first show that R(T) = 3. Clearly it is at most 3. Now suppose for the purpose of contradiction that  $a_1$  and  $a_2$  were spanned by two rank one slices  $b_1$  and  $b_2$ . It cannot be the case that both  $b_1$  and  $b_2$  have 0 in the upper left corner since then they would not span  $a_1$ . It cannot be the case that exactly one has a 0 in the upper left corner, because then that slice must equal  $a_2$  (which is not a rank one slice). So (after scaling) we must have

$$b_1 = \boxed{\begin{array}{c|c} 1 & x \\ y & xy \end{array}} \quad b_2 = \boxed{\begin{array}{c|c} 1 & s \\ t & st \end{array}}$$

from which we get the equation:

$$\begin{pmatrix} 1 & 1 \\ x & s \\ y & t \\ xy & st \end{pmatrix} \cdot M = \begin{pmatrix} 1 & 0 \\ 1 & 1 \\ 0 & 1 \\ 0 & 1 \end{pmatrix},$$

where M is a  $2 \times 2$  matrix. Thus M must be full rank, and so (y,t) = (xy, st). This equation leaves us with four possibilities: (1) y = 0 and t = 0; (2) y = 0 and s = 1; (3) x = 1 and t = 0; (4) x = 1 and s = 1. In case (1) we have (0, 0)M = (0, 1) from the third row of the linear system, a contradiction. In case (4) we have we have (1, 1)M = (1, 0)from the first row of the linear system and (1, 1)M = (1, 1) from the second row of the linear system, a contradiction. In case (2), we know that  $M = (1, 1; 0, t)^{-1} =$ (1, -1/t; 0, 1/t) from the first and third rows. But then (x, 1)M = (1, 1) implies x = 1, and this is a contradiction, since M has full rank (from the upper half of the linear system). Case (3) is symmetric; i.e., we know that  $M = (1, 1; y, 0)^{-1} = (0, 1/y; 1, 1/y)$ from the first and third rows. But then (1, s)M = (1, 1) implies s = 1, and this is a contradiction. We conclude that R(T) > 2.

Finally, the pseudorank of T is 2. Consider the slices

$$b_1 = \boxed{\begin{array}{c|c} 1 & 1 \\ 0 & 0 \end{array}} \quad b_2 = \boxed{\begin{array}{c|c} 1 & 2 \\ 1 & 2 \end{array}}$$

Note that  $a_1$  has the same support as  $b_1$  and  $a_2$  has the same support as  $b_2 - b_1$ .